

1 ■ Common formulae of information theory

2 Last compiled: August 14, 2007

3 Taken from T. Cover, J. Thomas, *Elements of Information Theory*, Wiley, 1991.

4 **Entropy.** $H(X) = - \sum_x p(x) \log p(x)$

$$H(X) \geq 0$$

$$H_b(X) = \log_b a H_a(X)$$

$$\log_a x = \log_a b \log_b x = \frac{\log_b x}{\log_b a}$$

5 $H(X) \leq \log |\mathcal{X}|$

$$d_\alpha \sum_x p(x)^\alpha|_{\alpha=1} = H(X)$$

$$(a^x)' = a^x \ln a, \quad (\log_a x)' = \frac{1}{x \ln a} = \frac{\log_a e}{x}$$

6 **Joint entropy.** $H(X, Y) = - \sum_{x,y} p(x, y) \log p(x, y)$

7 $H(X_1, \dots, X_n) \leq \sum_{i=1}^n H(X_i)$ eq. iff X_i ind.

8 **Conditional entropy.** $H(Y|X) = \sum_x H(Y|x)p(x) = - \sum_{x,y} p(x, y) \log p(y, x)$

9 $H(Y|X) \geq 0$ eq. if $Y = f(X)$, or iff $Y|x$ is deterministic for all $x \in \text{supp}(X)$

10 $H(Y|X) \leq H(Y)$ eq. iff X, Y ind.

11 **Relative entropy.** $D(p||q) = \sum_x p(x) \log \frac{p(x)}{q(x)}$

$$D(p||q) \geq 0$$

12 $D(p||u) = \log |\mathcal{X}| - H(X), \quad u(x) = |\mathcal{X}|^{-1}$

13 **Mutual information.** $I(X; Y) = \sum_{x,y} p(x, y) \log \frac{p(x, y)}{p(x)p(y)} = D(p(x, y)||p(x)p(y))$

$I(X; Y) \geq 0$ eq. iff X, Y ind.

$$I(X; Y) = I(Y; X)$$

$$I(X; X) = H(X)$$

14 $I(X; Y) = H(X) - H(X|Y) = H(Y) - H(Y|X)$

$$I(X; Y) = H(X) + H(Y) - H(X, Y)$$

$$\max I(X; Y) = \min(H(X), H(Y))$$

15 **Information metric.** $\Delta(X, Y) = H(X|Y) + H(Y|X)$

16 $\Delta(X, Y) = H(X) + H(Y) - 2I(X; Y) = H(X, Y) - I(X, Y)$

Chain rules.

$$H(X, Y) = H(X) + H(Y|X)$$

$$H(X, Y|Z) = H(X|Z) + H(Y|X, Z)$$

$$H(X_1, \dots, X_n) = \sum_{i=1}^n H(X_i|X_{i-1}, \dots, X_1)$$

$$I(X_1, \dots, X_n; Y) = \sum_{i=1}^n I(X_i; Y|X_{i-1}, X_{i-2}, \dots, X_1)$$

$$D(p(x, y)||q(x, y)) = D(p(x)||q(x)) + D(p(y|x)||q(y|x))$$

¹⁷ **Conditioning.**

$$\begin{aligned} D(p(y|x)||q(y|x)) &= \sum_x p(x) \sum_y p(y|x) \log \frac{p(y|x)}{q(y|x)} && \text{Conditional relative entropy} \\ I(X;Y|Z) &= H(X|Z) - H(X|Y,Z) && \text{Conditional mutual information} \end{aligned}$$

¹⁹ **Convexity properties.**

$$\begin{aligned} f(\lambda x_1 + (1-\lambda)x_2) &\leq \lambda f(x_1) + (1-\lambda)f(x_2) \text{ or } f''(x) \geq 0 && \text{Convex function} \\ E[f(X)] &\geq f(E[X]) && \text{Jensen's inequality} \\ \sum_{i=1}^n a_i \log \frac{a_i}{b_i} &\geq \left(\sum_{i=1}^n a_i \right) \log \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i} && \text{log sum inequality} \end{aligned}$$

²¹ **Differential entropy.** $h(X) = - \int p(x) \log p(x) dx$

$$\begin{aligned} h(X+a) &= h(X) && H(X^\Delta) + \log \Delta \rightarrow h(X) \\ h(AX) &= h(X) + \log |A| && D(P||Q) \geq 0 \\ h(X) &\leq \log \text{supp } X && D(P||Q) = \sup_\Delta D(P^\Delta||Q^\Delta) \\ h(X) &= \frac{1}{2} \log 2\pi e \sigma^2 \text{ when } X \sim N(m, \sigma^2) \end{aligned}$$

²³ **Correlations and causation.**

$$\begin{aligned} A \rightarrow B \leftarrow C &\quad I(A;C) = 0 \\ &\quad I(A;C|B) \neq 0 \text{ in general} \\ A \leftarrow B \rightarrow C \text{ or } &I(A;C) \neq 0 \text{ in general} \\ A \rightarrow B \rightarrow C &\quad I(A;C|B) = 0 \text{ (bottleneck)} \end{aligned}$$

²⁵ **Chains of random variables.**

$$\begin{aligned} X \rightarrow Y \rightarrow Z &\quad I(X;Y) \geq I(X;Z) \\ &\quad I(Z;Y) \geq I(Z;X) \text{ eq. iff } I(X;Y|Z) = 0 \\ I(X;Y) &\geq I(X;g(Y)) \\ I(X;Y|Z) &\leq I(X;Y) \\ X^{(n)} \rightarrow Y^{(k)} \rightarrow Z^{(m)} &\quad I(X;Z) \leq \log k \\ k < n, k < m &\quad I(X;Z) = 0 \text{ if } k = 1 \end{aligned}$$

²⁷ **Asymptotic equipartition theorem.**

²⁸ $X_1 X_2 \dots X_n \sim p(x)$ iid.

$$\begin{aligned} -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) &= -\frac{1}{n} \sum_{i=1}^n \log p(x_i) \\ &\rightarrow -E[\log p(x_i)] = H(X) \text{ (in probability)} \end{aligned}$$

$$\begin{aligned} A_\varepsilon^n &= \{x^n \in \mathcal{X}^n : 2^{-n(H+\varepsilon)} \leq p(x^n) \leq 2^{-n(H-\varepsilon)}\} && \text{Typical set} \\ -\frac{1}{n} \log p(x^n) &= H(X) \text{ (within } \varepsilon) \\ \Pr\{A_\varepsilon^n\} &> 1 - \varepsilon && \text{(from above result)} \\ |A_\varepsilon^n| &\leq 2^{n(H+\varepsilon)} \\ |A_\varepsilon^n| &\geq (1 - \varepsilon) 2^{n(H-\varepsilon)} && |A_\varepsilon^n| \doteq 2^{nH} \text{ (within } \varepsilon \text{ exponentially)} \\ p(x^n) &\doteq 2^{-nH} \end{aligned}$$

³⁰ **Method of types.**

$$\begin{aligned} X_1 X_2 \dots X_n, x_1 x_2 \dots x_n &= x^n = \mathbf{x} \in \mathcal{X}^n \\ P_{\mathbf{x}}(a) &= N(a|\mathbf{x})/n \\ \mathcal{P}_n &= \{P_{\mathbf{x}} : |\mathbf{x}| = n\} \\ T(P) &= \{\mathbf{x} \in X^n : P_{\mathbf{x}} = P\} \end{aligned}$$

$$\begin{aligned} \mathcal{X} &= \{a_1, a_2, \dots, a_{|\mathcal{X}|}\} \\ \sum_a P_{\mathbf{x}}(a) &= 1 \\ \text{Type} & \\ \text{Set of types } n & \\ \text{Type class} & \end{aligned}$$

$$|\mathcal{P}_n| \leq (n+1)^{|\mathcal{X}|}$$

$$\begin{aligned} |\mathcal{X}^n| &\sim |\mathcal{X}|^n \\ |\mathcal{P}_n| &\sim n^{|\mathcal{X}|} \\ X_1 X_2 \dots X_n &\sim \text{iid } Q(x) \\ \mathbf{x} \in T(Q) & \end{aligned}$$

$$\begin{aligned} \frac{1}{(n+1)^{|\mathcal{X}|}} 2^{nH(P)} &\leq |T(P)| \leq 2^{nH(P)} \\ \frac{1}{(n+1)^{|\mathcal{X}|}} 2^{-nD(P||Q)} &\leq Q^n(T(P)) \leq 2^{-nD(P||Q)} \end{aligned}$$

$$\begin{aligned} |T(P)| &\doteq 2^{nH(P)} & \text{Type class size} \\ Q^n(T(P)) &\doteq 2^{-nD(P||Q)} & \text{Type class prob.} \end{aligned}$$

$$\begin{aligned} \Pr\{D(P_{\mathbf{x}}||Q) > \varepsilon\} &\leq 2^{-n[\varepsilon - |\mathcal{X}| \frac{\log(n+1)}{n}]} \\ \Pr\{D(P_{\mathbf{x}}||Q) > \varepsilon\} &\leq n^{|\mathcal{X}|} 2^{-n\varepsilon} \sim 2^{-n\varepsilon} \end{aligned}$$

$$\begin{aligned} X_1 X_2 \dots X_n &\sim \text{iid } Q(x) & \text{WLLN} \\ D(P_{\mathbf{x}}||Q) &\rightarrow 0 \text{ (ip)} & \end{aligned}$$

$$\begin{aligned} Q^n(E) &\leq (n+1)^{|\mathcal{X}|} 2^{-nD(P^*||Q)} \\ P^* &= \arg \min_{P \in E} D(P||Q) \end{aligned}$$

$$\begin{aligned} X_1 X_2 \dots X_n &\sim \text{iid } Q(x) & \text{Sanov} \\ P &\subseteq \mathcal{P} & \end{aligned}$$

³² **Rate distortion theory.**

$$X^n \rightarrow f(X^n) \longrightarrow g(f(X^n)) \rightarrow \hat{X}^n$$

$$d : \mathcal{X} \times \hat{\mathcal{X}} \rightarrow \mathbb{R}^+ \quad d(x^n, \hat{x}^n) = \frac{1}{n} \sum_i d(x_i, \hat{x}_i) \quad \text{Distance, distortion}$$

$$d = \begin{cases} 0, & x = \hat{x} \\ 1, & x \neq \hat{x} \end{cases} \quad \max d(x, \hat{x}) < \infty \quad \text{Hamming distance}$$

$$R(D) = \min_{p(x|x) : E[d(X^n, \hat{X}^n)] \leq D} I(X; \hat{X}) \quad E[d(X, \hat{X})] = \Pr\{X \neq \hat{X}\} \quad f(X^n) \in \{1, 2, \dots, 2^{nR}\} \quad \text{Rate (\# bits needed)}$$

$$R(D) = \begin{cases} \frac{1}{2} \log \frac{\sigma^2}{D}, & 0 \leq D \leq \sigma^2 \\ 0, & D > \sigma^2 \end{cases} \quad \text{Gaussian channel}$$

³⁴ **Elements of probability theory.**

$$X_i \sim \text{iid}, \frac{1}{n} \sum_{i=1}^n X_i \xrightarrow{\text{ip}} E[X] \quad \text{WLLN}$$

$$Y = g(X), X = h(Y) = g^{-1}(Y) \quad f_Y(y) = f_X(h(y)) |h'(y)| = f_X(x) \left| \frac{\partial h(y)}{\partial y} \right|$$

$$Y = g(X) = \alpha X + \beta \quad \mathcal{N}(\mu, \sigma^2) \xrightarrow{g} \mathcal{N}(\alpha\mu + \beta, \alpha^2\sigma^2)$$

$$Y = X + Z \quad X \sim \mathcal{N}(x, \sigma_X^2)$$

$$Z \sim \mathcal{N}(z, \sigma_Z^2) \quad Y \sim \mathcal{N}(x+z, \sigma_X^2 + \sigma_Z^2)$$