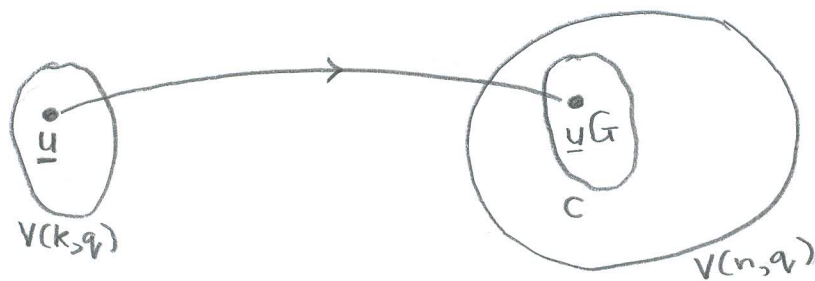


## Encoding with a linear code

Let  $C$  be a  $q$ -ary  $[n, k]$ -code.

Since  $M = q^k$ , we can communicate  $q^k$  messages — identify these with the  $q^k$  vectors of  $V(k, q)$ . If  $C$  has generator matrix  $G$ , then we encode these  $q^k$  message vectors with the map

$$V(k, q) \rightarrow V(n, q) : \underline{u} \mapsto \underline{u}G$$



Example Consider the binary  $[7, 4]$ -code  $C$

with generator matrix  $G = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$

The message vectors are the  $2^4 = 16$  vectors of  $V(4, 2)$ .

A message vector  $u_1 u_2 u_3 u_4$  is encoded by

$$[u_1 u_2 u_3 u_4] \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

$$= \left[ \underbrace{u_1 \quad u_2 \quad u_3 \quad u_4}_{\substack{\text{message} \\ \text{digits}}} \quad \underbrace{u_1+u_2+u_3 \quad u_2+u_3+u_4 \quad u_1+u_2+u_4}_{\substack{\text{check} \\ \text{digits}}} \right]$$

(added redundancy for protection against noise)

The  $2^4 = 16$  codewords are obtained by encoding the  $2^4 = 16$  message vectors; for example, the message vector 1011 is encoded to the codeword 1011000.

### Cosets

Let  $C$  be a  $q$ -ary  $[n, k]$ -code.

For any  $\underline{a} \in V(n, q)$ , the set

$$\underline{a} + C = \{ \underline{a} + \underline{x} : \underline{x} \in C \}$$

is called a coset of  $C$ .

Remember:

$V(n, q)$  is a group under vector addition and  $C$  is a subgroup of  $V(n, q)$ .

Lemma 3.3 If  $\underline{a}+C$  is a coset of  $C$

and  $\underline{b} \in \underline{a}+C$ , then  $\underline{b}+C = \underline{a}+C$ .

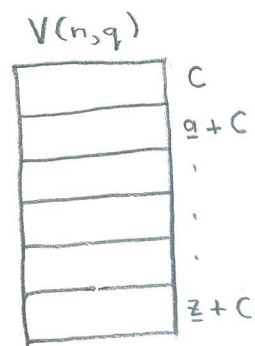
Theorem 3.4 Lagrange's Theorem

Suppose  $C$  is a  $q$ -ary  $[n, k]$ -code. Then

1. Every vector of  $V(n, q)$  belongs to some coset of  $C$ .

2. Every coset contains  $q^k$  vectors.

3. Any two cosets are disjoint.



Example: Let  $C$  be the binary  $[4, 2]$ -code with generator matrix  $G = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \end{bmatrix}$ .

Then  $C = \{0000, 1011, 0101, 1110\}$  and the

cosets of  $C$  are

$$0000 + C = C$$

$$1000 + C = \{1000, 0011, 1101, 0110\}$$

$$0100 + C = \{0100, 1111, 0001, 1010\}$$

$$0010 + C = \{0010, 1001, 0111, 1100\}$$

Note that, as representative for each coset, we have chosen a vector of minimum weight in the coset.

